Search, propagation, and learning in sequencing and scheduling problems

Mohamed Siala

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Context

Sequencing and Scheduling: the organization in time of of operations subject to capacity and resource constraints.



Context

PhD Context

- Combinatorial (optimization) problems
- Constraint satisfaction and optimization
- Laboratory: LAAS-CNRS, Toulouse
- Research Team: ROC (Operations Research, Combinatorial Optimization and Constraints)
- Supervision: Christian Artigues, and Emmanuel Hebrard
- Funding:



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Thesis overview

Constraint Programming: Search \oplus Propagation

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Thesis overview

Constraint Programming: Search \oplus Propagation \oplus Learning

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Thesis overview

Constraint Programming: Search \oplus Propagation \oplus Learning

All these aspects are important and must all be taken into account in order to design efficient solution methods

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Outline

Context

Background

- Case Study: The Car-Sequencing Problem
 Propagation
 - Learning
- 4 Learning in Disjunctive Scheduling
- 5 Conclusions & Perspectives

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A constraint is a finite relation

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A constraint is a finite relation

Definition

A constraint network (CN) is defined by a triplet $\mathcal{P} = (\mathcal{X}, \mathcal{D}, \mathcal{C})$ where

- $\mathcal{X} = [x_1, \dots, x_n]$: finite set of variables
- \mathcal{D} : a domain for \mathcal{X}
- C: finite set of constraints

A constraint is a finite relation

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- CSP is NP-Hard in general

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- Constraint Satisfaction Problem (CSP): deciding whether a constraint network has a solution or not
- CSP is NP-Hard in general
- Complete backtracking algorithms

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Search

• Search: decisions to explore the search tree

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Search

- Search: decisions to explore the search tree
- Search in CP= variable ordering + value ordering

Search

- Search: decisions to explore the search tree
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- Standard or customized

Search

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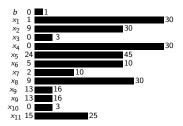
Propagation

- Propagation: inferences based on the current state
- Constraint \leftrightarrow propagator
- The level of pruning \leftrightarrow local consistency

Background

Learning

 $\begin{array}{l} x_1 + x_7 \geq 4 \wedge \\ x_2 + x_{10} \geq 11 \wedge \\ x_3 + x_9 = 16 \wedge \\ x_5 \geq x_8 + x_9 \wedge \\ b \leftrightarrow (x_9 - x_4 = 14) \wedge \\ b \rightarrow (x_6 \geq 7) \wedge \\ b \rightarrow (x_6 + x_7 \leq 9) \wedge \\ x_{11} \geq x_9 + x_{10} \end{array}$

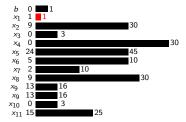


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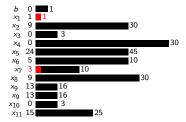


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$$[\![x_1=1]\!] \longrightarrow [\![x_7 \ge 3]\!]$$

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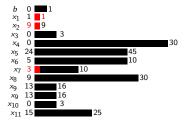


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$$\llbracket x_1 = 1 \rrbracket \longrightarrow \llbracket x_7 \ge 3 \rrbracket$$

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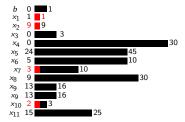
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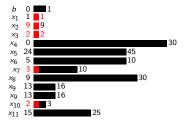
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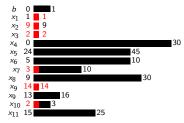
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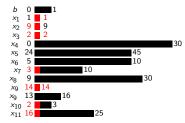
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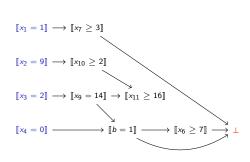
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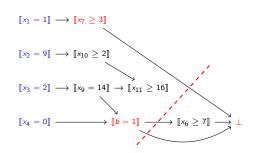
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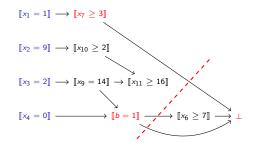


• Conflict analysis: $\llbracket b = 1 \rrbracket \land \llbracket x_7 \ge 3 \rrbracket \Rightarrow \bot$

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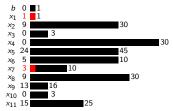


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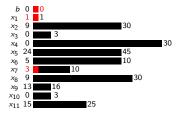


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- Propagate the learnt clause

$$\begin{array}{l} x_1 + x_7 \geq 4 \wedge \\ x_2 + x_{10} \geq 11 \wedge \\ x_3 + x_9 = 16 \wedge \\ x_5 \geq x_8 + x_9 \wedge \\ b \leftrightarrow (x_9 - x_4 = 14) \wedge \\ b \rightarrow (x_6 \geq 7) \wedge \\ b \rightarrow (x_6 + x_7 \leq 9) \wedge \\ x_{11} \geq x_9 + x_{10} \end{array}$$

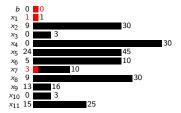


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- New clause: $\llbracket b \neq 0 \rrbracket \lor \llbracket x_7 \leq 2 \rrbracket$
- Backtrack to level 1
- Propagate the learnt clause
- Continue exploration

$$\begin{array}{l} x_1 + x_7 \geq 4 \wedge \\ x_2 + x_{10} \geq 11 \wedge \\ x_3 + x_9 = 16 \wedge \\ x_5 \geq x_8 + x_9 \wedge \\ b \leftrightarrow (x_9 - x_4 = 14) \wedge \\ b \rightarrow (x_6 \geq 7) \wedge \\ b \rightarrow (x_6 + x_7 \leq 9) \wedge \\ x_{11} \geq x_9 + x_{10} \end{array}$$



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Learning in CP

- $\bullet~{\rm Hybrid}~{\rm CP}/{\rm SAT}$
- Conflict Driven Clause Learning (CDCL) [Moskewicz et al., 2001]
- Based on the notion of explanation

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Summary of the thesis

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Modern CP-Solvers may not underestimate any of the three aspects: search, propagation, and learning

Contributions

• Propagation in a class of sequencing problems

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Modern CP-Solvers may not underestimate any of the three aspects: search, propagation, and learning

Contributions

- Propagation in a class of sequencing problems
- Search in car-sequencing

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- Propagation in a class of sequencing problems
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Contributions

- Propagation in a class of sequencing problems
- Search in car-sequencing
- Learning in car-sequencing
- Revisiting lazy generation

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Contributions

- Propagation in a class of sequencing problems
- Search in car-sequencing
- Learning in car-sequencing
- Revisiting lazy generation
- Learning in disjunctive scheduling

Outline



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Case Study: The Car-Sequencing Problem
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Car-Sequencing



- ROADEF'05 challenge [Solnon et al., 2008]
- RENAULT

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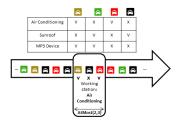
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• A class of vehicles is defined by a set of options

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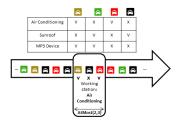
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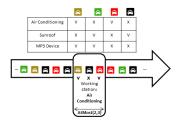
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- A class of vehicles is defined by a set of options
- Each class is associated to a demand
- Capacity constraints: no subsequence of size *q* may contain more than *p* vehicles requiring a given option
- Is there a sequence of cars satisfying both demand and capacity constraints?

Outline



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Case Study: The Car-Sequencing Problem
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Definition

 $\operatorname{ATMOSTSEQCARD}(\rho, q, d, [x_1, \dots, x_n]) \Leftrightarrow$

$$\bigwedge_{i=0}^{n-q} (\sum_{l=1}^{q} x_{i+l} \leq p) \land (\sum_{i=1}^{n} x_i = d)$$

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Definition

 $\operatorname{ATMOSTSEQCARD}(p,q,d,[x_1,\ldots,x_n]) \Leftrightarrow$

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Example ATMOSTSEQCARD($2, 5, 4, [x_1, \dots, x_9]$)

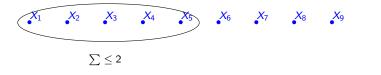
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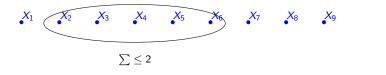
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Example ATMOSTSEQCARD($2, 5, 4, [x_1, \ldots, x_9]$)

$$\sum = 4$$

$$X_1 \quad X_2 \quad X_3 \quad X_4 \quad X_5 \quad X_6 \quad X_7 \quad X_8 \quad X_9$$

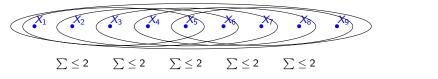
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Definition

 $\operatorname{ATMOSTSEQCARD}(p,q,d,[x_1,\ldots,x_n]) \Leftrightarrow$

$$\bigwedge_{i=0}^{n-q} \left(\sum_{l=1}^{q} x_{i+l} \leq p\right) \wedge \left(\sum_{i=1}^{n} x_{i} = d\right)$$

Example ATMOSTSEQCARD($2, 5, 4, [x_1, \dots, x_9]$)



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Example ATMOSTSEQCARD($2, 5, 4, [x_1, \ldots, x_9]$)

 X_1 X_2 X_3 X_4 X_5 X_6 X_7 X_8 X_9

- Car sequencing
- Crew-Rostering/Timetabling

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Arc Consistency on $\operatorname{AtMostSeqCard}$

Definition

A constraint C is Arc Consistent (AC) iff for every x in the scope of C, for every value $v \in D(x)$ there exists an assignment w in D satisfying C in which v is assigned to x

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$\bullet\ \operatorname{AtMostSeq} \oplus \operatorname{Cardinality}$ is not enough

ATMOSTSEQCARD as a particular case?

- COST-REGULAR: $O(2^q n)$ [van Hoeve et al., 2009]
- GEN-SEQUENCE: $O(n^3)$ [van Hoeve et al., 2009]
- GEN-SEQUENCE: $O(n^2.log(n))$ down a branch \oplus initial compilation of $O(q.n^2)$. [Maher et al., 2008].

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Arc Consistency on $\operatorname{AtMostSeqCard}$

An example with $\operatorname{ATMOSTSEQCARD}(4, 8, 12, [x_1, \dots, x_{22}])$. 0 0 1 0 1

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An example with ATMOSTSEQCARD(4, 8, 12, $[x_1, ..., x_{22}]$) . 0 0 1 0 1 . 0 0 1 0 1

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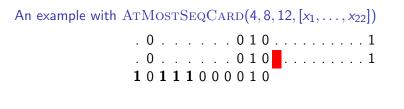
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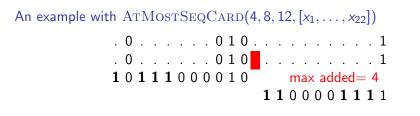
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Arc Consistency on ATMOSTSEQCARD



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Arc Consistency on $\operatorname{AtMostSeqCard}$

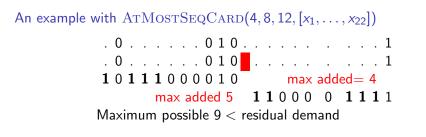


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Propagation

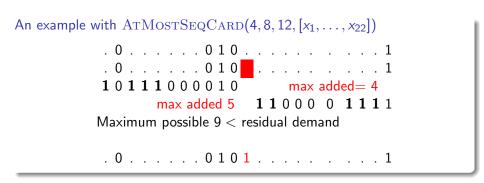
Arc Consistency on $\operatorname{AtMostSeqCard}$



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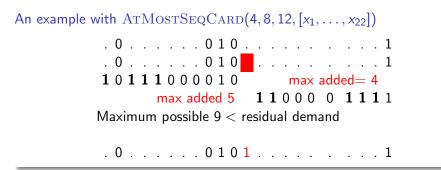
Propagation

Arc Consistency on $\operatorname{AtMostSeqCard}$



Propagation

Arc Consistency on $\operatorname{AtMostSeqCard}$



• Arc Consistency in O(n) time (optimal)

• Extremely efficient in practice (Car-Sequencing + Crew Rostering)

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Outline



Background

Case Study: The Car-Sequencing Problem
 Propagation

Learning

4 Learning in Disjunctive Scheduling

5 Conclusions & Perspectives

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Hybrid $\operatorname{CP}/\mathsf{SAT}$ Models

 \bullet Models based on $\operatorname{AtMostSeqCard}$

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- \bullet Models based on $\operatorname{AtMOStSeqCard}$
- We have to explain $\operatorname{AtMostSeqCard}$

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Explaining ATMOSTSEQCARD?

- Explain ATMOSTSEQ and CARDINALITY
- Explaining the extra filtering: consider the naive explanation, then try to reduce it.

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Experimental Results

- CP, SAT, Hybrid CP/SAT models
- Finding solutions quickly: Propagation is very important to find solutions quickly when they exist.
- For proving unsatisfiability: Clause learning is by far the most critical factor.

Outline

1 Context

2 Background

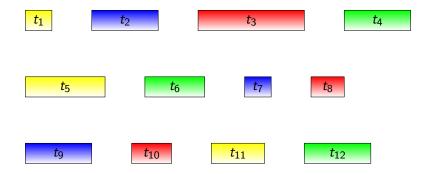
- Case Study: The Car-Sequencing Problem
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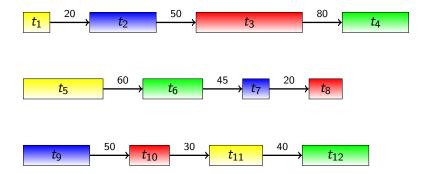
4 Learning in Disjunctive Scheduling

Conclusions & Perspectives

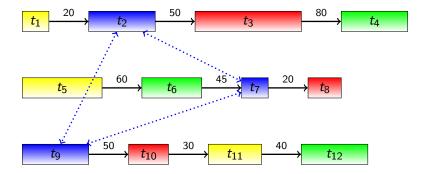
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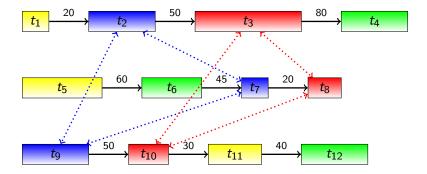




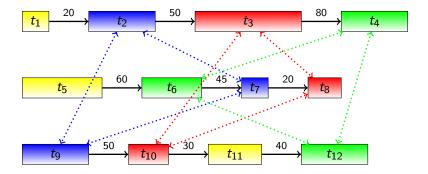
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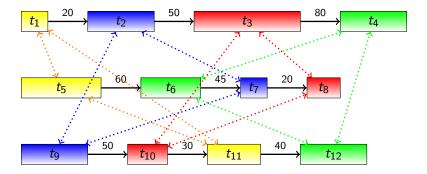
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Formulation

Unary Resource Constraint

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Formulation

Unary Resource Constraint

 \bullet Decomposition using the following $\ensuremath{\mathrm{DISJUNCTIVE}}$ constraints:

$$\delta_{kij} = \begin{cases} 0 \iff t_{ik} + p_{ik} \le t_{jk} \\ 1 \iff t_{jk} + p_{jk} \le t_{ik} \end{cases}$$
(1)

$\ensuremath{\mathrm{DISJUNCTIVE}}\xspace$ Learning

DISJUNCTIVE-based Learning

Conflict analysis in two phases:

- Standard 1-UIP cut
- Apply resolution for every bound literal until having a nogood with only reified Boolean variables

$\ensuremath{\mathrm{DISJUNCTIVE}}\xspace$ Learning

DISJUNCTIVE-based Learning

Conflict analysis in two phases:

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- \oplus Scheduling horizon does not manner in size
- \ominus Language of literals is restricted compared to standard approaches

$\ensuremath{\mathrm{DISJUNCTIVE}}\xspace$ Learning

DISJUNCTIVE-based Learning

Conflict analysis in two phases:

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new	old	new	old	new	old	new	old	new	old	new	old	new	old	
1305	1282	1613	1573	1514	1474	1543	1518	1561	1558	1573	1525	1508	1485	

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Modern constraint programming solvers may not underestimate any of these three aspects

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Modern constraint programming solvers may not underestimate any of these three aspects

Future Research

- (Car-Sequencing) Application to 'real' industrial situations?
- More extensions for ATMOSTSEQCARD?
- Hand crafted learning?

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Thank you for your attention!

Special thanks to my co-authors..

- Christian Artigues
- Emmanuel Hebrard
- Marie-Jose Huguet
- Valentin Mayer-Eichberger
- Nina Narodytska
- Thierry Petit
- Toby Walsh

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